

Computing Engineering

$$H_k = \frac{K_{e-d_k}}{m d_{k-1}}$$

$d_k$  maps free modules

so computing  $H_k$  is a linear alg problem.

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TOOL IS SMITH NORMAL FORM

which works over any PID

eg  $\mathbb{R}, \mathbb{Z}, (\mathbb{R}[\neq])$



$$\textcircled{4} \quad \mathbb{Z}^m / I_m(D) = \mathbb{Z}/d_1\mathbb{Z} \oplus \dots \oplus \mathbb{Z}/d_r\mathbb{Z} \oplus \mathbb{Z}^{n-r}$$

$$= \mathbb{Z}_{d_1} \oplus \dots \oplus \mathbb{Z}_{d_r} \oplus \mathbb{Z}^{n-r}$$

where  $\mathbb{Z}_{d_1} = \mathbb{Z}/d_1\mathbb{Z} = \mathbb{Z}/\mathbb{Z} = 0$

$$\Rightarrow \mathbb{Z}^m / I_m(M) = \text{Same.}$$

(5)  $A$  is  $m \times n$   $B$  is  $k \times m$   
 and  $BA = 0$

$\Rightarrow$  Given SNF of  $A$  and  $B$

$$\ker B / \text{Im } A = \mathbb{Z}^{m-r_B} / \text{Im } A$$

$$= \mathbb{Z}^{d_1} \oplus \dots \oplus \mathbb{Z}^{d_{l_A}} \oplus \mathbb{Z}^{m-r_B-l_A}$$

$d_L$  the diag on SNF(A).

$$\begin{array}{c}
 \underline{\underline{D_k^m}} \\
 C_{k-1} \xrightarrow{\partial_k} C_k \xrightarrow{\partial_{k+1}} C_{k+1} \\
 \partial_k \partial_{k+1} = 0 \text{ and } m_k = \dim \text{im}(C_k)
 \end{array}$$

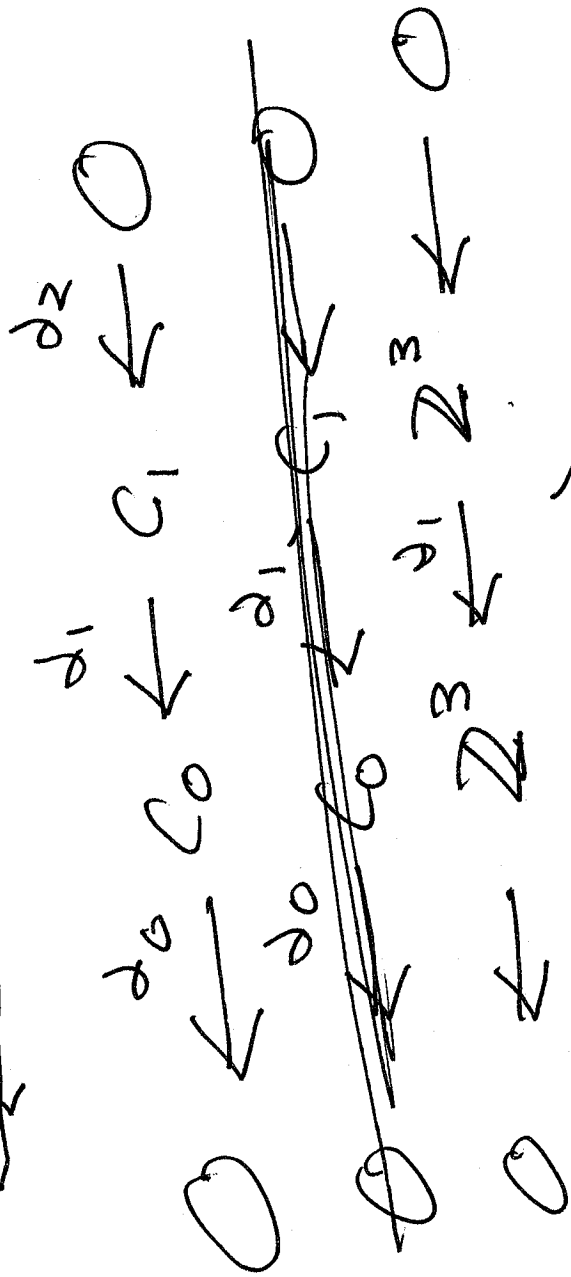
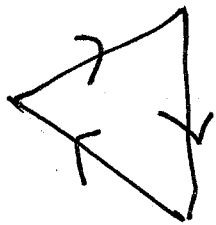
$D_k$  is SNF of  $\partial_k$ .

$$\begin{aligned}
 \Rightarrow H_k &= \mathbb{Z}^{d(k+1)} \oplus \cdots \oplus \mathbb{Z}^{d(k+1, l_{k+1})} \\
 &\oplus \mathbb{Z}^{m_k} = l(k) - l(k+1)
 \end{aligned}$$

Conclusion: Given the SNF of  $\partial_k$

$\Rightarrow H_k$  is computed.

Ex 11



$$\begin{bmatrix} -1 & 0 & 1 \\ 1 & -1 & 0 \\ 0 & 1 & -1 \end{bmatrix}$$

$$\begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 0 \end{bmatrix}$$

$d_1$

SNK

this page is purposely left blank. There is nothing missing.

$$m_2 = 0 \quad m_0 = 3 \quad m_1 = 3$$

$$l(1) = 2 \quad l(0) = \text{rank}(d_0) = 0$$

$$l(2) = 0$$

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$$m_0 - l(0) - l(1) = 3 - 0 - 2 = 1$$

$$H_0 = \mathbb{Z}$$

$$m_1 - l(1) - l(2) = 3 - 2 - 0 = 1$$

$$H_1 = \mathbb{Z}$$

no torsion since all  $d_i = 1$ .

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For how you can use computations

Plot form for SNF



~~CA~~ WARNING! SNF is slow

and expensive - poly in dimension

so efficient computation of  
homology is an active research area.

Now DATA  $\rightarrow$  Simplified  $\rightarrow$  Homology.  
 $\downarrow$  at scale  $\downarrow$  for scale  
complex  $\downarrow$   $\mathbb{Z}$   $\mathbb{Z}$

Persistent Homology

There are two basic methods to  
go from a finite data set to  
a simplicial complex

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$$X = \{x_0, \dots, x_k\} \subseteq \mathbb{R}^N \text{ (or metric space)}$$

distinct points

$$B_\varepsilon(x_i) = \{y \in \mathbb{R}^N : d(y, x_i) < \varepsilon\}$$

$$S \subseteq X$$

$$|B_\varepsilon(y)| = \{y \in S : d(y, S) < \varepsilon\}$$



Given  $X$  and  $\Sigma \geq 0$  be Čech complex  
 $\check{C}_\Sigma(X)$  is the abstract simplicial  
 complex with simplices

- (1) vertices are the points of  $X$
- (2) We have a simplex

$$\sum x_{a_1}, \dots, x_{a_l}$$

$$\Leftrightarrow \bigcap_{l=1}^L B_\Sigma(x_{a_l}) \neq \emptyset$$



Computationally expensive ~~to~~ ensure to  
~~decide~~ decide when balls intersect.

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The ~~altern~~ alternative more computable  
complex is Vietoris-Rips complex (VR)

given  $X$  and  $\epsilon > 0$ .  $VR_\epsilon(X)$



has simplices

(a) vertices points of  $X$

(b)  $\{x_{a_1}, \dots, x_{a_k}\} \Leftrightarrow$

$$d(x_{a_i}, x_{a_j}) < 2\epsilon$$

$\forall i, j$

(diameter of cluster of points)

Lemma 9

$$\cup C_{\Sigma}(\mathbb{X}) \subseteq \text{VR}_{\Sigma}(\mathbb{X}) \subseteq C_{2\Sigma}(\mathbb{X})$$