Approximate the lighting channel by a graph using LMA pulse |
locations as vertices in the graph. Edges are generated by
connecting closest neighboring vertices.

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William W. Hager |
hager@math.ufl.edu
http://www.math.ufl.edu/~hager
Department of Mathematics University of Florida
Gainesville, Florida 32611 USA 352-392-0281 $\times 244$

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LMA = rectangular matrix with 3 rows, each column gives the coordinates of an LMA pulse. The coordinates should be in meters.

EdgeList = Rectangular matrix with 2 rows, each column contains 2 integers which are the pulse number of two LMA pulses. These two pulses should be connected by a line segment and the union of these line segments approximate the lightning channel.
\% We remove pulses which are are isolated from all other pulses by a distance

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% given in the parameter node_connect. We grow the channel by first
% connecting together neighboring nodes to form channel segments.
% And then channel segments are connected together when
% they are near other channel segments. The parameter channel_connect
% gives the maximum separation distance that two channel segments can have
% and be connected together.
%
% We include an option to remove dangling nodes from the channel and to remove
% dangling channel segments. Set RemoveDangle = 1 to remove the dangling
% nodes and segments. A node n1 is dangling if it is connected to only one other
% node, say n2, and this other node n2 is connected to at least 2 other nodes:
% n1 (dangling node)
% ^
% |
% v
% n3 <--> n2 <--> n4
%
% A channel segment is considered dangling if the number of nodes in the
% channel is <= the value given in the variable DanglingChannel.
% Default parameter values:
% node_connect = 1000 (ignore nodes separated by 1 km from other nodes)
% channel_connect = 2000 (connect 2 channel segments if separation <= 2 km
% RemoveDangle = 1 (remove dangling nodes and channel segments)
% DanglingChannel = 15 (any channel segment with fewer than 15 nodes is
% ignored when RemoveDangle = 1)
function [EdgeList] = PulseGraph (LMA)
node_connect = 1000;
channel_connect = 2000;
RemoveDangle = 0;
DanglingChannel = 15;
nLMA = size (LMA, 2);
fprintf (1, 'Initial number LMA pulses: %i\n', nLMA) ;
nisolated = 0;
LMA_ones = ones (1, nLMA) ;
% initial pass, find all the isolated nodes
for j = 1:nLMA
    R = LMA - LMA (:, j)*LMA_ones ;
    Rnorm = sum (R.^2);
    Rnorm (j) = inf ;
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    [rmin k] = min (Rnorm);
    if ( rmin > node_connect^2 ) % skip this node
        LMA_ones (j) = 0 ;
        nisolated = nisolated + 1;
    end
end
fprintf (1, 'Number of isolated nodes: %i\n', nisolated) ;
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\% remove all isolated nodes

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LMA_keep = find (LMA_ones);
LMA = LMA (:, LMA_keep);
nLMA = size (LMA, 2);
fprintf (1, 'Number of LMA pulses (after removing isolates): %i\n', nLMA) ;
LMA_ones = ones (1, nLMA) ;
EdgeList = zeros (2, nLMA);
Sets = zeros (nLMA, 1) ;
Label = -ones (nLMA, 1) ; % initialize labels to -1
Next = zeros (nLMA, 1);
Prev = zeros (nLMA, 1);
nsets = 0;
nlabels =0;
nedges = 0;
```

\% pair up nodes that are adjacent to each other
for $\mathrm{j}=1$ : nLMA
if ( Label ( j ) $==-1$ )
R = LMA - LMA (:, j)*LMA_ones ;
Rnorm = sqrt (sum (R.^2));
Rnorm (j) = inf ;
[rmin k] = min (Rnorm);
if ( Label (k) ==-1) \% start a new set
nsets = nsets +1 ;
Sets (nsets) $=\mathrm{j}$;
Next (j) $=\mathrm{k}$;
$\operatorname{Prev}(\mathrm{j})=\mathrm{k}$;
Next (k) $=\mathrm{j}$;
$\operatorname{Prev}(k)=j$;
Label (j) = nsets ;
Label (k) = nsets ;
else $\quad \%$ add $j$ to set Label (k) ;
I = Prev (k) ;
$\operatorname{Prev}(\mathrm{j})=I$;
Next (I) $=\mathrm{j}$;

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        Next (j) = k;
        Prev (k) = j;
        Label (j) = Label (k) ;
    end
    nedges = nedges +1;
    EdgeList (1, nedges) = j ;
    EdgeList (2, nedges) = k ;
    end
end
fprintf (1, 'Initial number of edges: %i\n', nedges) ;
fprintf (1, 'Initial number of sets: %i\n', nsets) ;
ActiveSet = ones (nsets, 1) ;
nactive = nsets ;
AliveNodes = ones (nLMA, 1);
while ( nactive > 1) % add edges while no sets are isolated
% fprintf (1, 'begin loop, nactive: %i\n', nactive);
    i=0;
    while (i <= nsets )
        i= i + 1;
% fprintf (1, 'set # %i nsets: %i\n', i, nsets) ;
        if ( ActiveSet (i) ) % the set is not isolated from the rest
% fprintf (1,' active\n');
        j= Sets (i);
        jstart = j;
        LMA_ones = AliveNodes ;
        while (1)
            j = Next (j) ;
            LMA_ones(j) = 0 ;
            if ( }\textrm{j}==\textrm{j}\mathrm{ jtart )
                break;
            end
        end
        Comp = find (LMA_ones); % find the nodes in the complement
        lenComp = length (Comp);
% fprintf (1, 'size of complement: %i\n', lenComp) ;
        LMAComp = LMA (:, Comp) ;
        BestDist = inf ;
        ones_Comp = ones (1, lenComp) ;
        while (1)
        j = Next (j);
        R = LMAComp - LMA (:, j)*ones_Comp ;
        Rnorm = sqrt (sum (R.^2));
        [rmin k] = min (Rnorm);
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        if (rmin < BestDist )
            BestDist = rmin ;
            BestK = Comp (k) ;
            BestJ = j;
        end
        if ( j == jstart )
        break;
        end
end
% fprintf (1, 'dist to adjacent set: %e\n', BestDist) ;
if (BestDist <= channel_connect ) % connect two channel segments
    nactive = nactive - 1;
    nedges = nedges + 1;
    EdgeList (1, nedges) = BestJ ;
    EdgeList (2, nedges) = BestK ;
    I = Label (BestK) ;
    if (I< nsets ) % relabel the last set as I
        m = Sets (nsets);
        Sets (I) = m ;
        ActiveSet (I) = ActiveSet (nsets);
        n = m;
        while (1)
            n=Next (n);
            Label (n) = l;
            if ( }\textrm{n}==\textrm{m}\mathrm{ )
                break;
            end
        end
    end
    nsets = nsets - 1;
    % update the labels
    j = BestK ;
    while (1)
    j = Next (j) ;
        Label (j) = Label (Best) ;
        if ( }j== BestK
            break;
        end
    end
    % update the links
    m = Prev (BestK) ;
    n = Next (Best);
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                Next (BestJ) = BestK ;
            Prev (BestK) = BestJ ;
            Next (m)= n;
            Prev (n) = m;
            else % this set is separated from the rest
            fprintf(1,' kill\n');
            ActiveSet (i) = 0; % this set is dead
            nactive = nactive - 1;
            j = jstart ;
            while (1)
                j = Next (j) ;
                AliveNodes (j) = 0;
                if ( j == jstart )
                    break;
                    end
                    end
            end
        end
    end
end
% remove dangling nodes and their associated edge
if (RemoveDangle )
    degree = zeros (nLMA, 1);
    adj = zeros (nLMA, 1);
    edge = zeros (nLMA, 1);
    for j = 1:nedges
        k = EdgeList (1, j) ;
        I= EdgeList (2, j) ;
        adj (k) = I ;
        adj (I) = k;
        edge (k) = j;
        edge (l) = j;
        degree (k) = degree (k) + 1;
        degree (I) = degree (I) + 1 ;
    end
    J = find (degree == 1);
    EdgeFlag = ones (nedges, 1);
    ndangling = 0;
    for j = J'
        k = adj (j) ;
        if (degree (k) > 2 )
            degree (k) = degree (k) - 1;
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    EdgeFlag (edge (j)) = 0;
    ndangling = ndangling +1;
    end
end
fprintf (1, 'number of dangling nodes: %i\n', ndangling) ;
K = find (EdgeFlag);
EdgeList = EdgeList (:, K) ;
nedges = size (EdgeList, 2);
% remove dangling sets and their associated edge
ndrop = 0;
DropList = zeros (nsets, 1) ;
for i=1:nsets
    j = Sets (i) ;
    jstart = j;
    n = 0;
    while (1)
        j = Next (j) ;
        n=n+1;
        if ( }\textrm{j}== jstart
            break;
        end
    end
    if ( }\textrm{n}\mathrm{ < DanglingChannel ) % drop the set
        ndrop = ndrop + 1;
        DropList (ndrop) = i ;
    end
end
fprintf (1, 'number of dangling sets: %i\n', ndrop) ;
EdgeLabels = zeros (2, nedges);
for k = 1:nedges
    EdgeLabel (1, k) = Label (EdgeList (1, k)) ;
    EdgeLabel (2, k) = Label (EdgeList (2, k)) ;
end
EdgeFlag = ones (nedges, 1) ;
for i = 1:ndrop
    j = DropList (i) ; % drop this set
    J = find (EdgeLabel (1, :) == j) ;
    EdgeFlag (J) = 0;
    J = find (EdgeLabel (2, :) == j) ;
    EdgeFlag (J) = 0;
end
K = find (EdgeFlag);
EdgeList = EdgeList (:, K) ;
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    nedges = size (EdgeList, 2);
else
    EdgeList = EdgeList (:, 1:nedges) ;
end
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fprintf (1, 'final number of edges: \%i\n', nedges) ;
xmax = -inf ;
$x m i n=\inf$;
$y m a x=-i n f ;$
ymin $=\inf$;
LMA $=\mathrm{LMA} / 1000$;
plot the channel
figure ;
hold;
for $\mathrm{j}=1$ :nedges
$\mathrm{k}=$ EdgeList (1, j) ;
I = EdgeList ( $2, \mathrm{j}$ ) ;
x1 = LMA (1, k) ;
$x \max =\max (x \max , x 1)$;
$x \min =\min (x \min , x 1)$;
$\mathrm{x} 2=\mathrm{LMA}(1, \mathrm{I})$;
$x \max =\max (x \max , x 2)$;
$x \min =\min (x \min , x 2)$;
$\mathrm{X}=[\mathrm{x} 1 \times 2]$;
y1 = LMA (2, k) ;
$y \max =\max (y m a x, y 1)$;
$y \min =\min (y \min , y 1) ;$
$y 2=\operatorname{LMA}(2, I)$;
ymax $=\max (y m a x, y 2)$;
$y \min =\min (y \min , y 2) ;$
$\mathrm{Y}=[\mathrm{y} 1 \mathrm{y} 2]$;
plot (X, Y) ;
end
axis equal
axis ([xmin xmax ymin ymax]) ;
hold;
\% map back to original nodes
for $\mathrm{j}=1$ :nedges
EdgeList (1, j) = LMA_keep (EdgeList (1, j)) ;
EdgeList ( $2, \mathrm{j}$ ) $=$ LMA_keep (EdgeList $(2, \mathrm{j}))$;
end

